

Cartesian Model Structure

Knotting Problem

Many constructions in category theory produce data that only characterizes objects up to unique isomorphism.

One can make choices of representing but these choices will usually fail to be functorial.

Ex: $f: \text{Cat}(A, B)$ we obtain functors
 $\text{Mod}(A) \xrightarrow{\text{forget}} \text{Mod}(B) \xrightarrow{d^1} \text{Mod}(B/A)$
 d^1 is induced by $B \xrightarrow{\text{id}} B/A$
 d^1 is induced by $B \xrightarrow{\text{id}} B/A$

Now (B/A) of $\text{Set}^{A/B}$:
but d_1^1 and d_2^1 are only canonically isomorphic to the functor induced by c .

Classically one gets past this lack of functoriality by either of the following approaches:

(a) Redefining the notion of functor by saying we have a pseudofunctor
 $A \xrightarrow{F} \text{Cat}$ (so the unit and associativity conditions are functorial are relaxed to be local isoms)

(b) Constructing a dual fib.

where $r^*(\text{Set}) \cong \text{Mod}(B^{A/B})$
and give $M \in \text{Mod}$ and a map $(A) \xrightarrow{f} (B) \in A^*$ we have a map $M \xrightarrow{f^*} M \in A$ with $f^* = f$ and this f is determined by its universal property.
(Lif we chose another $f': M' \xrightarrow{f'} M$ with this property we obtain canonical equivalences $M' \cong M$.)

Approach (b) avoids making choices and instead emphasizes a property (such choices can be made). This makes this approach a bit easier to construct and a lot easier once we start trying to construct maps that respect our data.

The implication (a) \Rightarrow (b) is a right adjoint (or regular) which justifies (b).

(b) \Rightarrow (a) (using (a)) is essentially regarding a commutative diagram as a directed commutative diagram.

Let \mathcal{C} be Set^A w/ Joyal model str
 $\mathcal{C}^1 = \text{Set}^A / S$ w/ marked model str
 $\mathcal{C}^2 = \mathcal{C}^1$ w/ me. then mark every set $X^A \xrightarrow{p} S$ is cartesian
 $\mathcal{C}^3 = \text{Set}^A / S$ w/ Set^A marked str
 $\mathcal{C}_q = \text{Set}^A / \text{Set}^A$ w/ Set^A marked str

Then there is a chain of Quillen left adjoints

$$\mathcal{C} \xrightarrow{\text{id}} \mathcal{C}^1 \xrightarrow{\text{id}} \mathcal{C}^2 \xrightarrow{\text{id}} \mathcal{C}^3 \xrightarrow{\text{id}}$$

where \mathcal{C}^1 is cartesian
 \mathcal{C}^2 is cartesian
 \mathcal{C}^3 is cartesian

In categories it is even less natural to make choices since most categories arise from a multitude of possible models and/or from various universal properties.

Recollections (from Kashiwabara's talk)

Defn Let $p: X \rightarrow S$ be an inner fibration.
 $d^1: X \xrightarrow{\text{counit}} S$ is p -cart if the induced map $X_{/p} \xrightarrow{\sim} X_{/d^1} \times_S S_{/\text{pt}}$ is a trivial Kan fib.

Defn $p: X \rightarrow S$ is a set is

a Cartesian fibration if

- 1) p is an inner fib.
 - 2) $\forall \overset{\curvearrowleft}{X} \xrightarrow{p} S$ add every $x \in X$ satisfying
- $s.t. p(x) = \text{pt}$ \exists exact edge
 $\overset{\curvearrowleft}{X} \xrightarrow{S}$ satisfying $x \in X$ s.t.
 $p(x) = \text{pt}$.

Defn A Marked Set (X, E) is marked
a pair $X \in \text{Set}$ and $E \subseteq X$.
 $\mathbb{E} \in \text{Set}$. A map of such is a map
of sets sending marked edges to
marked edges. Denote this cat
by Set^* . (This is a closed, closed cartesian category
s.t. by (C1, C2) see [Kashiwabara](#))

Construction: The obvious forgetful functor
 $\text{Set}^* \rightarrow \text{Set}$ has a left adjoint
 $\text{Set} \xrightarrow{\text{forget}} \text{Set}^*$ (mark only deg edges)

and a right adjoint
 $\text{Set} \xleftarrow{\text{forget}} \text{Set}^*$ (mark all edges)

Both these functors preserve limits
and are cartesian monoidal.

These make Set^* into a
simplicially enriched (nonadditive) cartesian
category in two distinct ways:

$$\text{Set}^{\text{top}}((X, E), (Y, F)) = \text{Set}(X^{\text{op}}, Y)$$

$$\text{Set}^{\text{top}}((X, E), (Y, F)) = \text{Set}(X^{\text{op}}, Y^{\text{op}})$$

Conventions: Given a cartesian fib $p: X \rightarrow S$

Let $X^A = (X, E) \in \text{Set}^*$ where
 E is the set of p -cartesian edges
in X .

Ex: If $p: X \rightarrow S$ is invertible,
 X is an object that p is a
cartesian fib. and the marked
edges of X^A are precisely the
equivalences. (This is more general than previous)

Note: For Set^* , let $\text{Set}^{\text{top}} = \text{Set}^{\text{top}} \circ \text{Set}^*$.

Straightening for last term
 $X \xrightarrow{\text{forget}} \text{Set}$
set $S_{\text{pt}}(X) \in \text{Set}^{\text{top}}(\text{Set}^*, \text{Set}^{\text{top}})$
By $\text{forget}^{\text{top}}: (\text{Set}^{\text{top}})^{\text{op}} \times \text{Set}^{\text{top}}$
Note: $\text{Set}^{\text{top}} \cong \text{Set}^{\text{top}}$...
induces $\text{Set}^{\text{top}} \cong (\text{Set}^{\text{top}})^{\text{op}} \cong \text{Set}^{\text{top}}$

For $(X, E) \in \text{Set}^*$ we obtain

$S_{\text{pt}}(X) \in \text{Set}^{\text{top}}(\text{Set}^{\text{top}}, \text{Set}^{\text{top}})$

simplicial functor is as before
we just need to specify which edges

to mark in $S_{\text{pt}}(X)$ for Set^{top}

give $\beta \in S_{\text{pt}}(X)$, obtain $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$

$\beta = \beta^0 \xrightarrow{\text{id}} \beta^1$ and $\beta^0: X \rightarrow S$

$\beta^1: \overset{\curvearrowleft}{X} \rightarrow S$ defines

$\overset{\curvearrowleft}{X}$ is 1-simplicial

β^0 is 0-simplicial

Thin (with definitions to follow)

Let Set^* . Then \mathbb{I} is a left proper combinatorial model structure

- Combinatorics are precisely the Set^{top} (Set^{top} which are admissible)
- Weak equivalences are the contractible equivalences

Moreover with respect to the tensoring w/ Set^* via $\text{Set}^{\text{top}} \xrightarrow{\text{forget}} \text{Set}^*$ and how $\text{Set}^{\text{top}} \xrightarrow{\text{forget}} \text{Set}^*$ is a Set^* -model cat.

Similarly for $\text{Set}^{\text{top}} \xrightarrow{\text{forget}} \text{Set}^{\text{top}}$ and $\text{Set}^{\text{top}} \xrightarrow{\text{forget}} \text{Set}^{\text{top}}$ is a Set^{top} -model cat.

The fibrant objects of Set^* are precisely the objects of the form $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$ for some cartesian fib $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$.

Result $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$ is a cartesian even if $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$ is weakly cartesian

If $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$ is a cat even

if either of these conditions hold

(1) $\forall \text{Set}^*$

$\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$ is a cat even

(2) $\forall \text{Set}^*$ pk induces

a weak equivalence
of marked Kan Kancomplexes

Ex: If $\overset{\curvearrowleft}{Y} \xrightarrow{\beta} S$ cat fib
then $\text{Map}_{\text{Set}^*}((\beta), Y^A)$ is an object
and $\text{Map}_{\text{Set}^*}((\beta), Y^A)$ is the largest
Kan-subcomplex

Prop/Defn $S \in \text{Set}^*$ $\text{forget}^{\text{top}}(X, Y)$

f is a cat even if
either of the following equiv conditions
holds

1) $\forall Z \xrightarrow{\text{cat fib}}$
 $\text{Map}_{\text{Set}^*}(Y, Z) \xrightarrow{\text{id}} \text{Map}_{\text{Set}^*}(X, Z)$
is a cat even

2) $\forall Z \xrightarrow{\text{cat fib}}$
 $\text{Map}_{\text{Set}^*}(Y, Z) \xrightarrow{\text{id}} \text{Map}_{\text{Set}^*}(X, Z)$
is a homotopy even

Remark: If $X \xrightarrow{\beta} Y$ is a map
 $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$ is a cat even
if and only if β is a cat even

TFAE

- 1) f is a cat even
- 2) f has a strong higher inverse
i.e. $\text{Id}_{\text{Set}^*} \cong \text{Map}_{\text{Set}^*}(X, Y)$
endow over a contractible Kan complex
and similarly for f^{-1}
- 3) f induces categorical
equivalences otherwise

We mark these edges and also
mark any edge obtained by id_{Set^*}
along the composition of $(\text{Set}^{\text{top}})(\text{Set}^*)$,
that $\overset{\curvearrowleft}{X} \xrightarrow{\beta} S$ $\text{forget}^{\text{top}}(\text{Set}^{\text{top}}, \text{Set}^{\text{top}})$
is an Set^{top} equiv.